# 2006 Multi-Objective VRPTW Genetic Algorithm COSC 4F90

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- Vehicle Routing problem (VRP) is a specialization of the travelling salesman problem, making it NP complete.
  - This means there is (currently) no deterministic algorithm which can solve the problem in a reasonable (polynomial) amount of time.
- In VRPTW VRP with time windows we must also respect the opening and closing constraints of the customers we are servicing.
- The 2006 VRPTW paper, and this implementation, use a genetic algorithm in an attempt to find a 'good enough' solution to the VRPTW problem in a reasonable amount of time.

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- In essence, we are modelling evolution using computers to solve for a favourable solution using math. The advancement of this paper is the introduction of Pareto ranking.
  - Pareto ranking means we treat the problem as a multi-objective optimization problem instead of a single objective problem.
  - A solution is considered to be an improvement over another if it is as good or better in every objective that is part of the problem.
  - In this case, our objectives are reducing the total distance and total number of vehicles.

• Chromosome are just arrays.

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• Routes store the calculated best (valid) route from the chromosome, representing one vehicle.

```
struct chromosome {
    std::array<customerID_t, CUSTOMER_COUNT>
genes{};
};
struct route {
    ArrayList<customerID_t> customers;
    distance_t total_distance = 0;
};
```

### Individual and Population Representation

- Chromosomes are stored alongside all valid calculated routes within structures representing an individual in the population.
- the population stores only the set of individuals.

```
struct individual {
2
           chromosome c;
           ArrayList<route> routes{};
3
           distance_t total_routes_distance = 0;
4
           rank t rank = 0:
5
           fitness_t fitness = 0;
      }:
7
      struct population {
8
           ArrayList < individual > pops;
9
      };
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```

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### Genetic Algorithm Pseudo-code

```
read_instance_data(path);
1
      init_population();
2
      for (int i = 0; i < GENERATION_COUNT; i++){</pre>
3
           reconstruct_populations();
4
5
           calculatePopulationFitness();
6
           rankPopulation();
7
8
           population p;
9
           applyElitism(p, 1);
10
           while (p.size() < POPULATION_SIZE)</pre>
                applyTournamentAndOrCrossover(p);
12
13
           applyMutation(p);
14
           rebuild_population_chromosomes(p);
16
           current_population = p;
17
      }
18
```

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### **Tournament Selection**

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- The primary cause of my weird initial results.
- tournamentSelect() pseudo-code:

```
ArrayList <customerID_t> buffer;
fill_with_unique_individuals(buffer);
if (select(0,1) <= 0.8)
    return best_in(buffer);
else
    return random_from(buffer);
```

• applyTournamentAndOrCrossover pseudo-code:

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• applyCrossover(p1, p2) pseudo-code:

auto r1 = select\_rand\_route(p1.routes); auto r2 = select\_rand\_route(p2.routes, r1); auto c1, c2 = copy\_to\_children(p1, p2); // step a) remove p1's route from p2's (now c2) // remove p2's route from p1's (now c1) remove\_from(c2, r1); remove\_from(c1, r2); // insert r1 back into c2 at best and feasible point. insert\_to(c2, r1); // insert r2 back into c1 at best and feasible

- point.
- insert\_to(c1, r2);

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### Inserting back into a individual

#### • insert\_to(child, route) pseudo-code:

```
for (customer in route.customers){
    ArrayList<route> possibleRoutes;
    for (child_customer in child.routes.customers){
        auto r = insert_before_customer(customer, child_customer);
        if (feasible(r))
            possibleRoutes.push_back(r);
   // no feasible route found, we must make a new one
    if (possibleRoutes.empty())
        return new route(customer);
    else {
        route min:
        min.distance = double::max();
        for (route in possibleRoutes)
            if (route.distance < min.distance)
                min = route;
       // the actual implentation uses route_cache
       // to track the insertion location
        child.routes.insert_with_place(min);
```

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- In general, Pareto ranking performed as good to significantly better than weighted sum fitness.
- After fixing the index issue in tournament selection the R101 results got significantly better with Pareto ranking.
  - The results are significantly better than the Solomon best, which is suspicious...
- More on this later.

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### **R101 Pareto Results**

Solomon Best: Distance 1637.7, Routes 20



### C101 Pareto Results

Solomon Best: Distance 827.3, Routes 10



# R101 Weighted Sum Results

Solomon Best: Distance 1637.7, Routes 20



# C101 Weighted Sum Results

Solomon Best: Distance 827.3, Routes 10



- Because the c101 set contains values with a start time of 15 and an end time of 67 but service time of 90.00, it creates a situation where we cannot check the hard constraint of unloading before closing. All tests performed today are made by ignoring this issue.
- "lastDepartTime + record.service\_time > record.due"
- If an exception is made to allow routes in c101 to exist as single customer routes, the results become non-competitive with the Solomon best.
- Next slide contains a graph of r101 including the constraint above.

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# R101 Pareto Ranking

Solomon Best: Distance 1637.7, Routes 20



 Ombuki, Beatrice & Ross, Brian & Hanshar, Franklin. (2006). Multi-Objective Genetic Algorithms for Vehicle Routing Problem with Time Windows. Applied Intelligence. 24. 17-30. 10.1007/s10489-006-6926-z.

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